

Symbol Error Probability of Distributed-Alamouti Scheme in Wireless Relay Networks

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Abstract—In this paper, we analyze the maximum likelihood decoding performance of non-regenerative cooperation employing Alamouti scheme. Specifically, we derive two closed-form expressions for average symbol error probability (SEP) when the relays are located near by the source or destination. The analytical results are obtained as a single integral with finite limits and an integrand composed solely of trigonometric functions. Assessing the asymptotic (high signal-to-noise ratio) behavior of SEP formulas, we show that the distributed-Alamouti codes achieves a full diversity order. We also perform Monte-Carlo simulations to validate the analysis.

Index Terms—Alamouti, distributed space-time codes (DSTCs), non-regenerative relays, symbol error probability (SEP).

I. INTRODUCTION

A. Distributed-Alamouti Space-Time Coding Scheme

Alamouti scheme [1] has been considered as the only orthogonal space-time block code that can provide full rate and full diversity over complex constellation with symbol-wise maximum likelihood (ML) decoding complexity. Recently, there exists an extension of Alamouti scheme into cooperative/relay systems where relays simultaneously receive a noisy signal from the source and construct Alamouti space-time codes in a distributed fashion before relaying the signals to the destination [2]–[6].

B. Related Work

Distributed-Alamouti scheme dated back to the work in [2], [3], where two single-antenna relays assisted the source-destination communication by forming the Alamouti space-time block code in relay networks. The authors conjectured a diversity order of distributed-Alamouti scheme around one from their simulation. In [5], the average bit error rate (BER) was shown to be proportional to $\log(\text{SNR})/\text{SNR}^2$.

In [4], distributed-Alamouti system was created by using only one single-antenna relay in Protocol III, i.e., the relay and source construct a distributed-Alamouti space-time code and each terminal transmits each row of Alamouti code. Recently, exact closed-form expressions for pairwise error probability of this scheme has been analyzed in [6] where it has been shown that a full diversity order is achieved.

More recently, we generalized the construction of distributed space-time codes using amicable orthogonal design in relay networks [7]. It has been showed that our scheme can achieve full diversity order and single-symbol ML decoding complexity with the sacrifice of bandwidth efficiency.

C. Motivation and Results

Unlike most of previous works [2], [3], [5], [8], which focused on the average bit error rate (BER) of distributed-Alamouti space-time code for binary phase-shift keying (BPSK) modulation, we analyze the average symbol error probability (SEP) for M -ary phase-shift keying (M -PSK) by using the moment generating function (MGF) method [9], [10]. More specifically, we derive two closed-form expressions of average SEP for M -PSK modulation when the relays are close to either the destination or source. Since the asymptotic behavior of the MGF at large signal-to-noise (SNR) reveals a high-SNR slope of the SEP curve [9], [11], we show that distributed-Alamouti scheme achieves full diversity, i.e., a diversity order of two.

D. Organization and Notations

The paper is organized as follows. In the following section, we briefly review the cooperative system of two non-regenerative relays based on Alamouti scheme. We then derive two exact closed-form expressions of average SEP and diversity orders when two relays are closely located to both ends in Section III and Section IV, respectively. We show that two average SEP curves agree exactly with the Monte-Carlo simulation in Section V. Section VI provides the conclusion and future work.

Notation: Throughout the paper, we shall use the following notations. Vector is written as bold lower case letter and matrix is written as bold upper case letter. The superscripts $*$ and \dagger stand for the complex conjugate and transpose conjugate. \mathbf{I}_n represents the $n \times n$ identity matrix. $\|\mathbf{A}\|_F$ denotes Frobenius norm of the matrix \mathbf{A} and $|x|$ indicates the envelope of x . $\mathbb{E}_x\{\cdot\}$ is the expectation operator over the random variable x . A complex Gaussian distribution with mean μ and variance σ^2 is denoted by $\mathcal{CN}(\mu, \sigma^2)$. \log is the natural logarithm. $\Gamma(a, x)$ is the incomplete gamma function defined as $\Gamma(a, x) = \int_x^\infty t^{a-1} e^{-t} dt$ and $\mathcal{K}_0(\cdot)$ is the zeroth-order modified Bessel function of the second kind.

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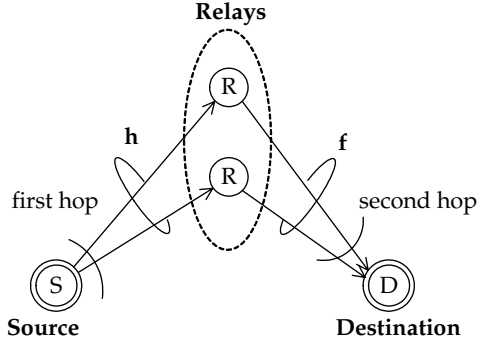


Fig. 1. Dual-hop non-regenerative relay channels.

II. SYSTEM MODEL

We consider a dual-hop relay channel, as shown in Fig. 1, where the channel remains constant for T_{coh} coherence time (an integer multiple of the dual-hop interval) and changes independently to a new value for each coherence time. All terminals are in half-duplex mode and, hence, transmission occurs over two time slots, each with the interval of two symbol periods.

In the first time slot, the source transmits two symbols $\mathbf{s} = [s_1 \ s_2]$, selected from M -PSK signal constellation \mathcal{S} , with average transmit power per symbol \mathcal{P}_s . During the first hop the received signals $\mathbf{y}_i = [y_i(1) \ y_i(2)]$ at i th relay is given by

$$\mathbf{y}_i = h_i \mathbf{s} + \mathbf{n}_{R_i} \quad (1)$$

where $y_i(j)$ is the j th symbol received at the i th relay, $h_i \sim \mathcal{CN}(0, \Omega_h)$ is the Rayleigh-fading channel coefficient for the source- i th relay link, \mathbf{n}_{R_i} is complex additive white Gaussian noise (AWGN) of variance N_0 , and $i = 1, 2$.

During the second time slot, the two relays construct Alamouti space-time scheme from the two received signals, and then retransmit a scaled version to the destination with the same power constraint as in the first hop, whereas the source remains silent. To simplify relaying operation, the relaying gain is determined only to satisfy the average power constraint with statistical channel state information (CSI) on \mathbf{h} (not its realizations) at the relay. With this *semi-blind* relaying, the output signal of the two relays are defined as

$$\mathbf{X} = \begin{bmatrix} x_1(1) & x_2(1) \\ x_1(2) & x_2(2) \end{bmatrix} = G \begin{bmatrix} y_1(1) & -y_2^*(2) \\ y_1(2) & y_2^*(1) \end{bmatrix} \quad (2)$$

where $x_i(j)$ is the j th symbol transmitted from the i th relay and G is the scalar relaying gain defined in the following. The received signal at the destination can be described as

$$r(j) = \sum_{i=1}^2 f_i x_i(j) + n_D(j) \quad (3)$$

where $f_i \sim \mathcal{CN}(0, \Omega_f)$ is the Rayleigh-fading channel coefficient for the i th relay-destination link and $n_D(j)$ is complex additive white Gaussian noise (AWGN) of variance N_0 . Note that all the random variable h_i and f_i , $i = 1, 2$ are statistically

independent and the variations in Ω_A , $A \in \{h, f\}$ capture the effect of distance-related path loss in each link. To constrain transmit power at the relay, we have

$$\mathbb{E} \left\{ \|\mathbf{X}\|_F^2 \right\} = 4G^2 (\Omega_h \mathcal{P}_s + N_0) = \mathbb{E} \left\{ \|\mathbf{s}\|_F^2 \right\} \quad (4)$$

yielding

$$G^2 = \frac{1}{2} \left(\Omega_h + \frac{1}{\text{SNR}} \right)^{-1} \quad (5)$$

where $\text{SNR} = \frac{\mathcal{P}_s}{N_0}$ is the common SNR of each link without fading [12]. The receive signals at the destination in (3) now can be equivalently described in the matrix form as

$$\mathbf{r} = \mathbf{H}\mathbf{s} + \mathbf{n} \quad (6)$$

where

$$\mathbf{r} = \begin{bmatrix} r(1) \\ r^*(2) \end{bmatrix},$$

$$\mathbf{H} = G \begin{bmatrix} h_1 f_1 & -h_2^* f_2 \\ h_2 f_2^* & h_1^* f_1^* \end{bmatrix},$$

which is complex orthogonal, i.e.,

$$\mathbf{H}^\dagger \mathbf{H} = G^2 \sum_{i=1}^2 |h_i|^2 |f_i|^2 \mathbf{I}_2,$$

and

$$\mathbf{n} = G \begin{bmatrix} f_1 n_{R_1}(1) & -f_2 n_{R_2}^*(2) \\ f_1^* n_{R_1}^*(2) & f_2^* n_{R_2}(1) \end{bmatrix} + \begin{bmatrix} n_D(1) \\ n_D^*(2) \end{bmatrix}$$

which is zero-mean AWGN of the variance $N_0 \left(G^2 \sum_{i=1}^2 |f_i|^2 + 1 \right)$. It is easy to see that the ML decoding of the symbol vector \mathbf{s} turns out very simple as two symbols in \mathbf{s} are independently decomposed from one another. Therefore, the instantaneous receive SNR is readily written as

$$\gamma = \text{SNR} \frac{G^2 \sum_{i=1}^2 |h_i|^2 |f_i|^2}{G^2 \sum_{i=1}^2 |f_i|^2 + 1} \quad (7)$$

To examine the statistical characteristic of γ given in (7) we consider two special cases. If the relays are much closer to the destination than the source, then we may have $G^2 \sum_{i=1}^2 |f_i|^2 \gg 1$. In this special case, we express γ in the following

$$\gamma = \text{SNR} \frac{\sum_{i=1}^2 |h_i|^2 |f_i|^2}{\sum_{i=1}^2 |f_i|^2} \quad (8)$$

On the other hand, if the relays are near the source, then the following expression may hold $G^2 \sum_{i=1}^2 |f_i|^2 \ll 1$. Therefore, the instantaneous receive SNR γ is as follows

$$\gamma = \text{SNR} G^2 \sum_{i=1}^2 |h_i|^2 |f_i|^2 \quad (9)$$

III. CLOSED-FORM EXPRESSION OF THE AVERAGE SEP AND DIVERSITY ORDER

In this section, on account of the statistical behavior of the instantaneous receive SNR for two special cases as shown in previous section, we now derive the close-form expression of the average SEP and then deduce the diversity order of distributed-Alamouti scheme in such cases.

A. When the relays are close to the destination

With the instantaneous receive SNR in (8), the MGF of γ , averaged over channel coefficients, is given by

$$\phi_\gamma(\nu) \triangleq \mathbb{E}_\gamma \{ \exp(-\nu\gamma) \} = \mathbb{E}_{h_i, f_i} \{ \exp(-\nu\gamma) \} \quad (10)$$

Since $A_i \sim \mathcal{CN}(0, \Omega_A)$, $A \in \{h, f\}$ and $i = 1, 2$, it is obvious that $|A_i|^2$ obeys an exponential distribution with hazard rate $1/\Omega_A$. The probability density function (p.d.f.) of $|A_i|^2$ can be written as

$$p_{|A_i|^2}(x) = \frac{1}{\Omega_A} \exp(-x/\Omega_A) \quad (11)$$

Since h_i and f_i are statistically independent with each other, the MGF of γ in (10) can be written as

$$\begin{aligned} \phi_\gamma(\nu) &= \mathbb{E}_{f_i} \{ \mathbb{E}_{h_i} \{ \exp(-\nu\gamma) \} \} \\ &= \mathbb{E}_{f_i} \left\{ \prod_{i=1}^2 \left(1 + \text{SNR}\nu \frac{|f_i|^2}{|f_1|^2 + |f_2|^2} \right)^{-1} \right\} \\ &= \mathbb{E}_z \left\{ \left(1 + \xi \frac{z}{z+1} \right)^{-1} \left(1 + \xi \frac{1}{z+1} \right)^{-1} \right\} \\ &\stackrel{(a)}{=} \int_0^\infty [(1 + (1 + \xi)z)(1 + \xi + z)]^{-1} dz \\ &= \frac{2 \log(1 + \xi)}{\xi(2 + \xi)} \end{aligned} \quad (12)$$

where $z = \frac{|f_1|^2}{|f_2|^2}$, $\xi = \Omega_h \text{SNR}\nu$, and (a) follows immediately from Theorem 2 in Appendix. Using the well-known MGF approach [10], [13] along with (12), we obtain the average SEP of the distributed-Alamouti scheme with M -PSK in relay channels as¹

$$\begin{aligned} P_e &= \frac{1}{\pi} \int_0^{\pi - \frac{\pi}{M}} \phi_\gamma \left(\frac{g_{\text{MPSK}}}{\sin^2 \theta} \right) d\theta \\ &= \frac{1}{\pi} \int_0^{\pi - \frac{\pi}{M}} \frac{2 \sin^4 \theta \log \left(1 + \frac{\psi}{\sin^2 \theta} \right)}{\psi (2 \sin^2 \theta + \psi)} d\theta \end{aligned} \quad (13)$$

where $\psi = \Omega_h \text{SNR} g_{\text{MPSK}}$ and $g_{\text{MPSK}} = \sin^2(\pi/M)$.

B. When the relays are close to the source

In this section, we will consider the case when the two relays are close to the source. Following the same steps as in

¹The result can be applied to other binary and M -ary signals in a straightforward way (see, e.g., [13]).

Section III-A and from (9), the MGF of γ can be described as

$$\phi_\gamma(\nu) = \mathbb{E}_{h_i, f_i} \left\{ \prod_{i=1}^2 \exp \left(-G^2 \text{SNR}\nu |h_i|^2 |f_i|^2 \right) \right\} \quad (14)$$

$$= \left[\int_0^\infty \exp(-G^2 \text{SNR}\nu t) p_T(t) dt \right]^2 \quad (15)$$

$$= \left[\int_0^\infty \frac{2}{\Omega} \exp(-G^2 \text{SNR}\nu t) \mathcal{K}_0 \left(2\sqrt{\frac{t}{\Omega}} \right) dt \right]^2 \quad (16)$$

$$= \left[\lambda \exp(\lambda) \Gamma(0, \lambda) \right]^2 \quad (17)$$

where $T = |h_i|^2 |f_i|^2$, $\Omega = \Omega_h \Omega_f$, $\lambda = (G^2 \text{SNR}\nu \Omega)^{-1}$, (16) follows immediately from Theorem 3 in Appendix, and (17) can be obtained from the change of variable $v = G^2 \text{SNR}\nu t$ along with [14, eq. (8.353.4)]. Therefore, similarly as in the previous subsection, the SEP is given by

$$P_e = \frac{1}{\pi} \int_0^{\pi - \frac{\pi}{M}} \left[\tau \sin^2 \theta \exp(\tau \sin^2 \theta) \Gamma(0, \tau \sin^2 \theta) \right]^2 d\theta \quad (18)$$

where $\tau = (G^2 \text{SNR}\nu \Omega g_{\text{MPSK}})^{-1}$.

C. High-SNR Characteristic: Diversity Order

We now assess the effect of cooperative diversity on the SEP behavior in a high-SNR regime. The diversity impact of non-regenerative cooperation on a high-SNR slope of the SEP curve can be quantified by the following theorem.

Theorem 1 (Achievable Diversity Order): The non-regenerative cooperation of distributed-Alamouti scheme provides maximum diversity order, i.e.,

$$D \triangleq \lim_{\text{SNR} \rightarrow \infty} \frac{-\log P_e}{\log(\text{SNR})} = 2. \quad (19)$$

Proof: See Appendix C ■

IV. NUMERICAL RESULTS

In this section, we validate our analysis by comparing with Monte-Carlo simulation. In the following numerical examples, we consider the non-regenerative relay protocol employing Alamouti code as in Section II. We assume collinear geometry for locations of three communicating terminals, as shown in Fig. 2. The path-loss of each link follows an exponential-decay model: if the distance between the source and destination is equal to d , then $\Omega_0 \propto d^{-\alpha}$ where the exponent $\alpha = 4$ corresponding to a typical non line-of-sight propagation. Then, $\Omega_h = \epsilon^{-\alpha} \Omega_0$ and $\Omega_f = (1 - \epsilon)^{-\alpha} \Omega_0$.

Fig. 3 and Fig. 4 show the SEP of QPSK versus SNR when the relay approaches the destination with $\epsilon = 0.7$ and $\epsilon = 0.8$ and the source with $\epsilon = 0.2$ and $\epsilon = 0.3$, respectively. As seen from both figures, analytical and simulated SEP curves match exactly. Observe that the PEP slopes for $\epsilon = 0.7$ and $\epsilon = 0.8$ are identical at the high SNR regime, as speculated in Theorem 1. The same observation can be obtained for $\epsilon = 0.2$ and $\epsilon = 0.3$. For comparison, two examples demonstrate the

performance is decreased when the relay approaches both end, e.g., the SEP for $\epsilon = 0.3$ is slightly less than that for $\epsilon = 0.2$ and a 3dB-gain in SEP performance can be obtained with $\epsilon = 0.7$ compared to the case with $\epsilon = 0.8$.

APPENDIX

A. Auxiliary Results

The following lemmas will be useful in the paper

Lemma 1: Let $a > 0$ be a finite constant and

$$f(x) = \frac{2 \log(1+ax)}{ax(2+ax)}, \quad x > 0 \quad (20)$$

We have

$$f_{x\uparrow} \triangleq \lim_{x \rightarrow \infty} \frac{-\log f(x)}{\log x} = 2 \quad (21)$$

Proof: It follows immediately from (20) and (21) that

$$\begin{aligned} f_{x\uparrow} &= \underbrace{\lim_{x \rightarrow \infty} \frac{-\log \log(1+ax)}{\log x}}_{\stackrel{(b)}{=} 0} + \underbrace{\lim_{x \rightarrow \infty} \frac{\log(ax)}{\log x}}_{\stackrel{(c)}{=} 1} \\ &\quad + \underbrace{\lim_{x \rightarrow \infty} \frac{\log(1+\frac{ax}{2})}{\log x}}_{\stackrel{(d)}{=} 1} = 2 \end{aligned} \quad (22)$$

where (b), (c), and (d) follow immediately by applying l'Hôpital rule. ■

Lemma 2: Let

$$g(\zeta) = [\zeta \exp(\zeta) \Gamma(0, \zeta)]^2, \quad \zeta > 0 \quad (23)$$

Let $\zeta = \alpha \frac{\beta x + 1}{x^2}$ and $\alpha, \beta > 0$ be finite constants. We have

$$g_{x\uparrow} \triangleq \lim_{x \rightarrow \infty} \frac{-\log g(x)}{\log x} = 2 \quad (24)$$

Proof: Substituting $\zeta = \alpha \frac{\beta x + 1}{x^2}$ into (23), it follows immediately from (24) that

$$\begin{aligned} g_{x\uparrow} &= -2 \left[\lim_{x \rightarrow \infty} \frac{-\log \left(\frac{\beta x + 1}{x^2} \right)}{\log x} + \lim_{x \rightarrow \infty} \frac{\beta x + 1}{x^2 \log x} \right. \\ &\quad \left. + \lim_{x \rightarrow \infty} \frac{\Gamma \left(0, \alpha \frac{\beta x + 1}{x^2} \right)}{\log x} \right] = -2 \left[-2 + \underbrace{\lim_{x \rightarrow \infty} \frac{\log(1+\beta x)}{\log x}}_{\stackrel{(e)}{=} 1} \right. \\ &\quad \left. + \underbrace{\lim_{x \rightarrow \infty} \frac{\beta x + 1}{x^2 \log x}}_{\stackrel{(f)}{=} 0} + \underbrace{\lim_{x \rightarrow \infty} \frac{\Gamma \left(0, \alpha \frac{\beta x + 1}{x^2} \right)}{\log x}}_{\stackrel{(g)}{=} 0} \right] = 2 \end{aligned} \quad (25)$$

where (e), (f) follow immediately by l'Hôpital rule and (g) follows from the Laguerre² polynomial series representation of incomplete gamma function [14] together with l'Hôpital rule. ■

²The incomplete gamma function can be described as $\Gamma(0, x) = e^{-x} \sum_{n=0}^{\infty} \frac{L_n(x)}{n+1}$ where $L_n(x) = \sum_{m=0}^n (-1)^m (n!x^m) / (m!(n-m)!m!)$ is the Laguerre polynomial of order n .

B. A Ratio and Product Distribution

Theorem 2 (Ratio Distribution): Let

$$X \sim \Upsilon(1/\Omega)$$

$$Y \sim \Upsilon(1/\Omega)$$

be statistically independent and identically distributed (i.i.d.) exponential r.v.'s. Suppose the ratio Z of the form

$$Z = \frac{X}{Y} \quad (26)$$

Then, we obtain the p.d.f. of random variable Z as

$$p_Z(z) = (z+1)^{-2} \quad (27)$$

Proof: Note that

$$\begin{aligned} p_Z(z) &= \int_0^{\infty} y p_{XY}(yz, y) dy \\ &= \int_0^{\infty} \frac{y}{\Omega^2} \exp\left(-y \frac{z+1}{\Omega}\right) dy \\ &= (z+1)^{-2} \end{aligned} \quad (28)$$

Theorem 3 (Product Distribution): Let

$$X \sim \Upsilon(1/\Omega_x)$$

$$Y \sim \Upsilon(1/\Omega_y)$$

be statistically independent and not necessarily identically distributed (i.n.i.d.) exponential r.v.'s. Suppose the product T of the form

$$T = XY \quad (29)$$

Then, we have

$$p_T(t) = \frac{2}{\Omega_x \Omega_y} \mathcal{K}_0 \left(2 \sqrt{\frac{t}{\Omega_x \Omega_y}} \right) \quad (30)$$

where $\mathcal{K}_0(\cdot)$ is the zeroth-order modified Bessel function of the second kind.

Proof: Note that

$$\begin{aligned} F_T(t) &= \Pr\{XY \leq t\} \\ &= \mathbb{E}_X \{F_{T|X}(t)\} \\ &= \mathbb{E}_X \left\{ 1 - \exp\left[-\frac{t}{x\Omega_y}\right] \right\} \\ &= 1 - \frac{1}{\Omega_x} \int_0^{\infty} \exp\left[-\frac{t}{x\Omega_y} - \frac{x}{\Omega_x}\right] dx \end{aligned} \quad (31)$$

The p.d.f. of T follows immediately from differentiating (31) with respect to t .

$$\begin{aligned} p_T(t) &= \frac{1}{\Omega_x \Omega_y} \int_0^{\infty} \frac{1}{x} \exp\left[-\frac{t}{x\Omega_y} - \frac{x}{\Omega_x}\right] dx \\ &= \frac{2}{\Omega_x \Omega_y} \mathcal{K}_0 \left(2 \sqrt{\frac{t}{\Omega_x \Omega_y}} \right) \end{aligned} \quad (32)$$

where the last equality follows from the change of variable $u = \frac{x}{\Omega_x}$ along with [14, eq. (8.432.6)] as desired. ■

C. Proof of Theorem 1

Since the asymptotic behavior of the MGF $\phi_\gamma(\nu)$ at large SNR reveals a high-SNR slope of the SEP curve, we have [11]

$$D = \lim_{\text{SNR} \rightarrow \infty} \frac{-\log \phi_\gamma(g_{\text{MPSK}})}{\log(\text{SNR})}. \quad (33)$$

Hence, it follows immediately from (12) and Lemma 1 with $a = \Omega_h g_{\text{MPSK}}$ that $D = 2$ when the relays are close to the destination. Also, from (17) and Lemma 2 with $\alpha = \frac{2}{\Omega_h \Omega_f g_{\text{MPSK}}}$ and $\beta = \Omega_h$, we can obtain $D = 2$ as in (19) when the relays are much close to the source.

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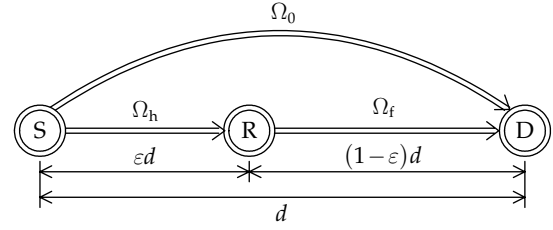


Fig. 2. Collinear topology with an exponential-decay path loss model where $\Omega_0 \propto d^{-\alpha}$, $\Omega_h = \varepsilon^{-\alpha} \Omega_0$, and $\Omega_f = (1-\varepsilon)^{-\alpha} \Omega_0$ with $\alpha = 4$.

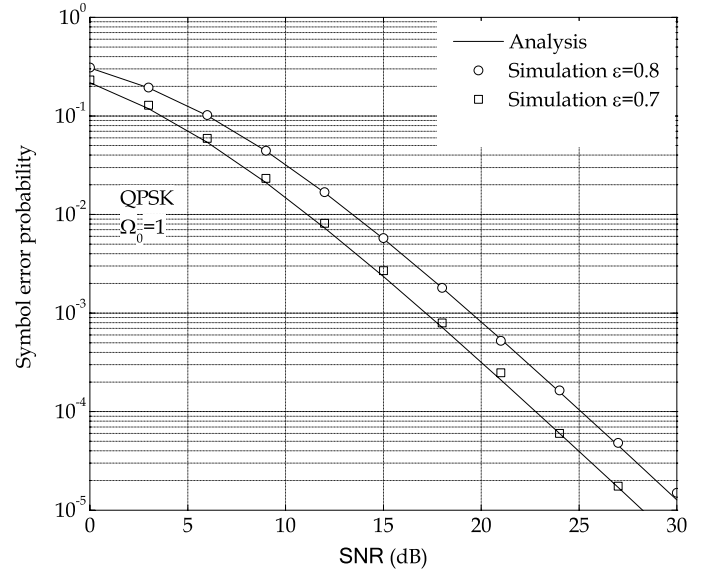


Fig. 3. Symbol error probability of QPSK versus SNR in non-regenerative relay channels employing Alamouti scheme when $\varepsilon = 0.7$ and $\varepsilon = 0.8$ (close the destination). $\Omega_0 = 1$.

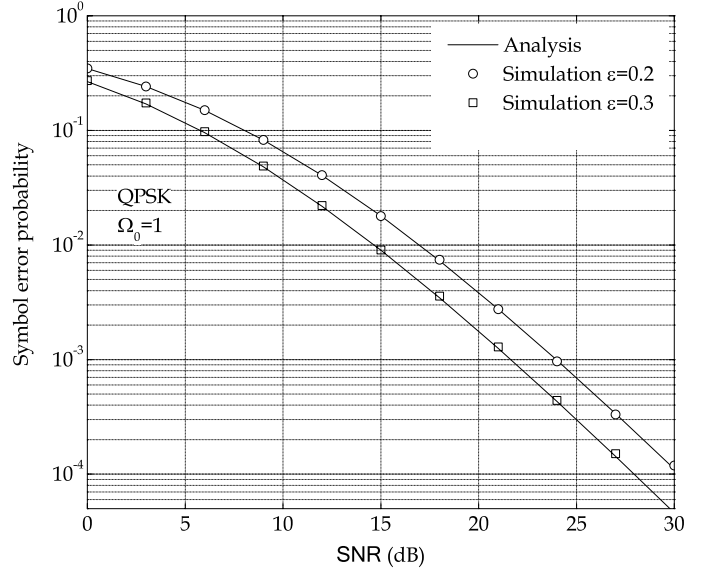


Fig. 4. Symbol error probability of QPSK versus SNR in non-regenerative relay channels employing Alamouti scheme when $\varepsilon = 0.2$ and $\varepsilon = 0.3$ (close the source). $\Omega_0 = 1$.